

## ON THE NUMBER OF GRAPHS WITHOUT 4-CYCLES

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The number  $F_n$  of  $n$ -vertex graphs lacking a 4-cycle is shown to satisfy  $\frac{1}{2}n^{3/2} \leq \log F_n \leq 1.08192n^{3/2}$  for large  $n$ . This represents an improvement of the previously known upper and lower bounds and resolves a long-standing question posed by Erdős. The result is obtained by using an argument previously used by the authors to bound the numbers of lattices on  $n$  elements.

### Introduction

In this paper we consider the number  $F_n$  of  $n$ -vertex labelled graphs without loops or multiple edges in which there is no four vertex cycle. We first give a construction which produces  $(\sqrt{2})^{n^{3/2}}$  graphs of the desired type. It is not hard to find an upper bound of the form  $n^{cn^{3/2}}$  ( $c$  a constant) for  $F_n$ .

The question as to which of these two expressions is closer to  $F_n$  has been open. Erdős some years ago posed the question: is the logarithm of  $F_n$  asymptotic to (or bounded in) both directions by a constant multiple of  $n^{3/2}$ , or by a constant multiple of  $(\log n)n^{3/2}$ ? In this note we resolve this question, showing that the logarithm of  $F_n$  lies between two multiples of  $F_n$ , as follows:

$$\frac{1}{2}n^{3/2} \leq \log_2 F_n + o(n^{3/2}) \leq (1.081919 \dots)n^{3/2}.$$

The problem is closely related to that of enumerating lattices on  $n$  elements. One important class of lattices are those in which chains have no more than four elements; that is, in which there are (in addition to the minimum and maximum elements) only atoms and co-atoms. Kleitman and Winston [2] recently found an upper bound of the form  $2^{cn^{3/2}}$  ( $c$  a constant) for the number of such lattices.

An undirected bipartite graph in which there are no 4-cycles can be taken as the graph of the covering relation between atoms and co-atoms of such a lattice. The bound on the number of these lattices could then be applied iteratively to

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yield a bound of the desired form for general 4-cycle-free graphs. However, a better upper bound can be obtained by applying the argument directly to graphs.

### The lower bound

We use a property of projective planes admitting polarities to construct graphs of the desired type. A projective plane of order  $q$  has a (point, line) incidence matrix  $A = (a_{ij})$  in which  $a_{ij}$  is zero if point  $i$  does not lie on line  $j$ , and  $a_{ij}$  is one if point  $i$  does lie on line  $j$ .  $A$  is a  $(q^2 + q + 1)$  by  $(q^2 + q + 1)$  matrix, and every row and column sum is  $q + 1$ . Furthermore, for every pair of distinct rows  $i$  and  $j$  of  $A$ , there is exactly one  $k$  with  $a_{ik} = a_{jk} = 1$ ; the dual property holds for columns.

A polarity of a projective plane is an incidence-preserving one-to-one map of order 2 between points and lines and between lines and points. If a plane admits a polarity, we can assume that point  $i$  is mapped into line  $i$ . This means that the incidence matrix  $A$  of such a plane is symmetric.

We construct from such a matrix  $A$  a labelled graph  $G$  with  $n = q^2 + q + 1$  vertices by including an edge between vertex  $i$  and vertex  $j$  if and only if  $i \neq j$  and  $a_{ij} = a_{ji} = 1$ . Clearly  $G$  has no 4-cycles.

Furthermore, any subgraph of  $G$  lacks 4-cycles. There are at least  $q(q^2 + q + 1)$  ones off the main diagonal of  $A$ , so there are at least  $\frac{1}{2}n(\sqrt{n - \frac{1}{4}} - \frac{1}{2})$  edges in  $G$ . Thus

$$\log_2 F_n \geq \frac{1}{2}n^{3/2} + c(n^{3/2}) \quad (1)$$

if there is a projective plane of order  $q$  admitting a polarity, with  $n = q^2 + q + 1$ .

Finally note that there is a Desarguan projective plane of every prime power order. Dembowski [1, p. 157] has shown that every Desarguan plane admits a polarity. For large  $n$ , the difference between  $n$  and the nearest prime power is small compared to  $n$ , so (1) holds for all  $n$ .

### The upper bound

Let  $d_i$  be the degree of the  $i$ th vertex of an  $n$ -vertex graph  $G$  lacking 4-cycles. Then  $\sum_{i=1}^n \binom{d_i}{2}$  counts the number of pairs of vertices both connected to a third vertex. Thus

$$\sum_{i=1}^n \binom{d_i}{2} \leq \binom{n}{2}; \quad (2)$$

otherwise some pair of vertices would both be connected to two other vertices, giving a 4-cycle. This leads to

$$\frac{1}{n} \sum_{i=1}^n (d_i^2 - d_i) \leq n - 1, \quad (3)$$

and therefore to

$$\left(\frac{1}{n} \sum_{i=1}^n d_i\right)^2 - \left(\frac{1}{n} \sum_{i=1}^n d_i\right) \leq n-1. \quad (4)$$

Solving the quadratic equation gives

$$\frac{1}{n} \sum_{i=1}^n d_i \leq \sqrt{n-\frac{3}{4}} + \frac{1}{2}. \quad (5)$$

The minimum degree  $d (= \min d_i)$  in an  $n$ -vertex graph without 4-cycles is therefore at most  $\lceil \sqrt{n} \rceil$ ; the greatest integer in the right hand side of (5) for large  $n$ .

Now note that  $F_n$  is at most  $n$  times the number of  $n$ -vertex graphs lacking 4-cycles in which the  $n$ th vertex is of minimum degree  $d$ . There are at most

$$\binom{n-1}{d} \leq \binom{n-1}{\lceil \sqrt{n} \rceil}$$

ways to add an  $n$ th vertex of degree  $d$  to an  $(n-1)$ -vertex graph lacking 4-cycles; thus

$$F_n \leq n F_{n-1} \binom{n-1}{\lceil \sqrt{n} \rceil} \leq F_{n-1} n^{\sqrt{n}}, \quad (6)$$

which gives the bound

$$\log_2 F_n \leq \frac{2}{3}(\log_2 n - \frac{2}{3})n^{3/2} + o(n^{3/2}) \quad (7)$$

when applied inductively.

We now demonstrate how to improve this upper bound so that the  $\log_2 n$  term on the right hand side of (7) is replaced by a constant. We will prove:

**Theorem.** *The number  $F_n$  of graphs with  $n$  vertices in which there is no 4-cycle satisfies*

$$\log_2 F_n \leq \alpha n^{3/2} + o(n^{3/2}),$$

where

$$\alpha = \frac{2}{3} \max_{0 \leq x \leq 1} \left( \frac{x^2 \log_2 x^2 - (1-x^2) \log_2(1-x^2)}{x} \right) n \sim 1.08192$$

We preface the proof with a discussion of the problem and a description of the argument.

We will obtain the upper bound by providing a construction procedure to add a vertex to an  $(n-1)$ -vertex graph lacking 4-cycles so that the resulting graph also lacks 4-cycles. Each  $n$ -vertex graph lacking 4-cycles can be produced by applying this procedure to some  $(n-1)$ -vertex graph lacking 4-cycles.

We will show that the factor of  $\binom{n-1}{\lceil \sqrt{n} \rceil}$  appearing in (6) can be replaced by a

factor that is closer to  $\binom{n/d}{d}$  for some number  $d$ ; this factor attains its maximum when  $d \sim \sqrt{n}/2.03$  and leads to the bound in the theorem.

The factor  $\binom{n-1}{n-1}$  was introduced in the inequality (6) above as a rough bound on the number of ways in which we could choose connections for a new minimum degree vertex. However, every time we choose a vertex with which the  $n$ th vertex (the vertex being added) is to share an edge, certain other vertices become ineligible to share an edge with the  $n$ th vertex because such an edge would complete a 4-cycle. This cuts down on the number of choices that can be made. In fact, it will turn out that after connecting an appropriately small number of vertices to the  $n$ th vertex, we can insure that there are at most  $n/d$  ( $d$  the minimum degree of the graph) vertices remaining which can share an edge with the  $n$ th vertex without completing a 4-cycle. This will provide the desired bound.

**Proof of the Theorem.** Let  $G_0$  be the  $(n-1)$ -vertex graph  $G$  with an isolated  $n$ th vertex added. We proceed in  $d$  steps, adding an edge at each step. At the first step all vertices except the  $n$ th are eligible to be connected (i.e. to share an edge with) the  $n$ th vertex.

At each step, form an ordered list of the eligible vertices as follows: After  $(i-1)$  vertices have been put on the list, determine which one of the remaining eligible vertices is connected in  $G$  by a path of length 2 to the greatest number of eligible vertices not yet on the list. (Ties can be broken arbitrarily.) That vertex will be  $i$ th on the list.

We then choose an eligible vertex from the list and add the edge between the  $n$ th vertex and the chosen vertex. All vertices which came before the chosen vertex are made ineligible, as are all vertices connected to the chosen vertex by a path of length 2. We then go to the next step, repeating until we either reach a dead end or form a graph in which the  $n$ th vertex has degree  $d$ .

Note that any  $n$ -vertex graph lacking 4-cycles (in which the  $n$ th vertex has minimum degree) can be formed from some  $(n-1)$ -vertex graph lacking 4-cycles with this procedure. The procedure provides a framework for producing these graphs; it does not limit the graphs that can be produced.

Suppose that there are  $s$  eligible vertices at some step and one vertex is connected at that step to the  $i$ th vertex on the list. We want to bound the number of eligible vertices left on the list after this choice. Let  $\bar{s} = s - i + 1$ ; there are  $\bar{s}$  vertices after (and including) the chosen vertex on the list. Let  $q_m$  be the number of these  $\bar{s}$  vertices with which the  $m$ th vertex of  $G$  shares an edge. Note that

$$\sum_{m=1}^{n-1} q_m \geq d\bar{s} \tag{8}$$

since all vertices have degree at least  $d$ . There are  $\sum_{m=1}^{n-1} \binom{\bar{s}}{2}$  pairs of the  $\bar{s}$  vertices connected by paths of length 2. By again making use of the fact that the average of a square is greater than the square of its average we may bound this number of

pairs according to

$$\sum_{m=1}^{n-1} \binom{q_m}{2} \geq \frac{d\bar{s}(d\bar{s}-n+1)}{2(n-1)}. \quad (9)$$

The chosen vertex is connected by a path of length 2 to the most others below it on the list. Therefore it is connected to at least

$$(2/\bar{s}) \sum_{m=1}^{n-1} \binom{q_m}{2} \geq \frac{d(d\bar{s}-n+1)}{(n-1)} \quad (10)$$

vertices further down the list by a path of length 2.

We can now bound the number of eligibles remaining on the list. There are at most

$$\bar{s} - \frac{d(d\bar{s}-n+1)}{(n-1)} \leq \bar{s} \left(1 - \frac{d^2}{n}\right) + d \quad (11)$$

eligibles remaining so that  $\bar{s} - n/d$  after one iteration becomes reduced to  $(\bar{s} - n/d)(1 - d^2/n)$ . Applying (11) repeatedly from the first step where  $\bar{s} \leq n-1$ , we see that after  $z$  steps there are at most  $n(1 - d^2/n)^z + n/d$  and hence at most  $\exp(-zd^2/n) + n/d$  remaining eligible vertices. Therefore if  $zn/d^2 \geq \log n$  the number of eligibles from which choices can be made must be at most  $n/d$ .

(6) can therefore be altered to read

$$F_n \leq nF_{n-1} \binom{n}{z} \binom{n/d}{d-z}, \quad (12)$$

$$z = \lceil (n/d^2 \log n) \rceil \quad (13)$$

where  $\binom{n}{z}$  overestimates the number of choices available in the early steps and  $\binom{n/d}{d-z}$  is obtained by assuming free choice after there are only  $n/d$  eligible vertices left.

To prove the theorem it suffices to show that

$$\log_2 \binom{n}{z} \binom{n/d}{d-z} < \frac{3}{2} \alpha \sqrt{n}, \quad (14)$$

(where  $\alpha$  is given in the statement of the theorem). Note that if

$$d < \frac{\frac{3}{2} \alpha \sqrt{n}}{\log_2 n} \quad \text{then} \quad \log \binom{n}{d} < \frac{3}{2} \alpha \sqrt{n},$$

so we need only consider values of  $d$  obeying

$$d \geq \frac{\frac{3}{2} \alpha \sqrt{n}}{\log_2 n}.$$

This means we can take  $z$  to be at most the order of  $(\log n)^3$ , so the  $\binom{n}{z}$  factor is negligible.

Using Sterling's formula, we have

$$\log \binom{n/d}{d-z} < \log_2 \binom{n/d}{d} < -d \log_2 \left( \frac{d^2}{n} \right) - \left( \frac{n}{d} - d \right) \log_2 \left( 1 - \frac{d^2}{n} \right). \quad (15)$$

Let  $x^2 = d^2/n$ . Then

$$\log_2 \binom{n/d}{d} < \sqrt{n} \frac{H(x^2)}{x}, \quad (16)$$

where  $H(y) = -y \log_2 y - (1-y^2) \log_2 (1-y^2)$  is the entropy function.

The right hand side of (16) attains a maximum of about 1.62288 when  $x \approx 0.49143$ . Thus

$$\log_2 \binom{n/d}{d} < \frac{3}{2}(1.081919 \dots) \sqrt{n}, \quad (17)$$

proving the theorem.

The number of bipartite graph of order  $n$  is of the form  $c^{n^2}$ , so the number of graphs without odd cycles of any given size is of order  $c^{n^2}$ . The number of graphs lacking 6-cycles, or both 4 and 6-cycles, has not been investigated. The problem may well be amenable to the type of argument used here.

## References

- [1] P. Dembowski, *Finite Geometries* (Springer-Verlag, New York, 1968).
- [2] D.J. Kleitman and K.J. Winston, The asymptotic number of lattices, *Ann. Discrete Math.* 6 (1980) 243-249.